

## Appendix A: Missing Definitions, Lemmas and Proofs

### A.1. Proof of Theorem 3

*Proof:* The proof of the running time is trivial. Because  $\pi^f$  takes  $k$  rounds to select  $k$  items, and the running time of each round is bounded by  $\frac{n}{k} \log \frac{1}{\epsilon}$  which is the size of the random set  $H$ . Thus, the total running time is bounded by  $O(k \times \frac{n}{k} \log \frac{1}{\epsilon}) = O(n \log \frac{1}{\epsilon})$ .

We next prove the approximation ratio of  $\pi^f$ . Let  $\phi'$  denote the worst-case realization for  $\pi^f$ , i.e.,  $\phi' = \arg \min_{\phi} \mathbb{E}[f(E(\pi^f, \phi), \phi)]$ . Let random variables  $\Psi'_t$  denote the partial realization after selecting the first  $t$  items  $\mathbf{S}_t$ , which is also a random variable, conditioned on  $\phi'$ . We can represent the expected worst-case utility  $\tilde{f}_{wc}(\pi^f)$  of  $\pi^f$  as follows:

$$\begin{aligned} \tilde{f}_{wc}(\pi^f) &= \mathbb{E}_{\Pi}[f(E(\pi^f, \phi'), \phi')] \\ &= \sum_{t \in [k]} \{ \mathbb{E}_{\Pi}[f(E(\pi_t^f, \phi'), \phi')] - \mathbb{E}_{\Pi}[f(E(\pi_{t-1}^f, \phi'), \phi')] \} \\ &= \sum_{t \in [k]} \{ \mathbb{E}_{\mathbf{S}_t, \Psi'_t}[f(\mathbf{S}_t, \Psi'_t)] - \mathbb{E}_{\mathbf{S}_{t-1}, \Psi'_{t-1}}[f(\mathbf{S}_{t-1}, \Psi'_{t-1})] \} \end{aligned} \quad (33)$$

To prove this theorem, it suffices to show that for all  $t \in [k]$ ,

$$\mathbb{E}_{\mathbf{S}_t, \Psi'_t}[f(\mathbf{S}_t, \Psi'_t)] - \mathbb{E}_{\mathbf{S}_{t-1}, \Psi'_{t-1}}[f(\mathbf{S}_{t-1}, \Psi'_{t-1})] \geq (1 - \epsilon) \frac{f_{wc}(\pi_{wc}^*) - \mathbb{E}_{\mathbf{S}_{t-1}, \Psi'_{t-1}}[f(\mathbf{S}_{t-1}, \Psi'_{t-1})]}{k} \quad (34)$$

This is because by induction, (34) and (33) imply that  $\tilde{f}_{wc}(\pi^f) \geq (1 - (1 - \frac{1-\epsilon}{k})^k) f_{wc}(\pi_{wc}^*) \geq (1 - e^{-(1-\epsilon)}) f_{wc}(\pi_{wc}^*) \geq (1 - 1/e - \epsilon) f_{wc}(\pi_{wc}^*)$ .

The rest of the proof is devoted to proving (34). Consider a given partial realization  $\psi'_{t-1}$  after selecting the first  $t-1$  items  $S_{t-1}$ . Let  $W \in \arg \max_{S \subseteq E: |S|=k} \sum_{e \in S} f_{wc}(e | \psi'_{t-1})$  denote the top  $k$  items that have the largest worst-case marginal utility on top of  $\psi'_{t-1}$ . Recall that  $H$  is a random set of size  $\frac{n}{k} \log \frac{1}{\epsilon}$  sampled uniformly at random from  $E$ , we first show that the probability that  $H \cap W = \emptyset$  is upper bounded by  $e^{-|H| \frac{k}{n}}$ .

$$\Pr[H \cap W = \emptyset] \leq \left(1 - \frac{|W|}{|E|}\right)^{|H|} = \left(1 - \frac{k}{n}\right)^{|H|} \leq e^{-|H| \frac{k}{n}}$$

It follows that  $\Pr[H \cap W \neq \emptyset] \geq 1 - e^{-|H| \frac{k}{n}}$ . Because we assume  $|H| = \frac{n}{k} \log \frac{1}{\epsilon}$ , it follows that

$$\Pr[H \cap W \neq \emptyset] \geq 1 - e^{-|H| \frac{k}{n}} \geq 1 - e^{-\frac{n}{k} \log \frac{1}{\epsilon} \frac{k}{n}} \geq 1 - \epsilon \quad (35)$$

Now we are ready to prove (34). Again, consider a fixed partial realization  $\psi'_{t-1}$  after selecting the first  $t-1$  items  $S_{t-1}$ , we have

$$\begin{aligned} &\mathbb{E}_{\mathbf{S}_t, \Psi'_t}[f(\mathbf{S}_t, \Psi'_t) | S_{t-1}, \psi'_{t-1}] - f(S_{t-1}, \psi'_{t-1}) \\ &= \mathbb{E}_{\mathbf{e}'_t}[f(S_{t-1} \cup \{\mathbf{e}'_t\}, \psi'_{t-1} \cup \{(\mathbf{e}'_t, \phi'(\mathbf{e}'_t))\}) - f(S_{t-1}, \psi'_{t-1}) | S_{t-1}, \psi'_{t-1}] \\ &\geq \mathbb{E}_{\mathbf{e}'_t}[f_{wc}(\mathbf{e}'_t | \psi'_{t-1}) | S_{t-1}, \psi'_{t-1}] \geq \mathbb{E}_H[\max_{e \in H} f_{wc}(e | \psi'_{t-1}) | S_{t-1}, \psi'_{t-1}] \\ &\geq \Pr[H \cap W \neq \emptyset] \frac{\sum_{e \in H} f_{wc}(e | \psi'_{t-1})}{k} \geq (1 - \epsilon) \frac{\sum_{e \in W} f_{wc}(e | \psi'_{t-1})}{k} \geq (1 - \epsilon) \frac{\sum_{i \in [k]} f_{wc}(e_i^* | \psi'_{t-1})}{k} \\ &\geq (1 - \epsilon) \frac{\sum_{i \in [k]} f_{wc}(e_i^* | \psi'_{t-1} \cup \psi_{i-1}^*)}{k} = (1 - \epsilon) \frac{f(S_{t-1} \cup S_k^*, \psi'_{t-1} \cup \psi_{k-1}^*) - f(S_{t-1}, \psi'_{t-1})}{k} \\ &\geq (1 - \epsilon) \frac{f(S_k^*, \psi_k^*) - f(S_{t-1}, \psi'_{t-1})}{k} \\ &= (1 - \epsilon) \frac{f_{wc}(\pi_{wc}^*) - f(S_{t-1}, \psi'_{t-1})}{k} \end{aligned} \quad (36)$$

The third inequality is due to each item of  $W$  has equal probability of being included in  $H \cap W$ . The fourth inequality is due to (35). The fifth inequality is due to  $W \in \arg \max_{S \subseteq E: |S|=k} \sum_{e \in S} f_{wc}(e | \psi'_{t-1})$ . The sixth inequality is due to (5). Taking the expectation of (36) over  $(\mathbf{S}_{t-1}, \Psi'_t)$ , we have (34).  $\square$

## A.2. Proof of Proposition 1

*Proof:* Consider any two partial realizations  $\psi$  and  $\psi'$  such that  $\psi \subseteq \psi'$  and any item  $e \in E \setminus \text{dom}(\psi')$ , we have

$$\begin{aligned} f_{wc}(e | \psi) &= \min_{o \in O(e, \psi)} f(\text{dom}(\psi) \cup \{e\}, \psi \cup \{(e, o)\}) - f(\text{dom}(\psi), \psi) \\ &= \min_{o \in O(e, \psi)} 1 - p_H(\mathcal{H}(\psi \cup \{(e, o)\})) - (1 - p_H(\mathcal{H}(\psi))) \\ &= \min_{o \in O(e, \psi)} p_H(\mathcal{H}(\psi)) - p_H(\mathcal{H}(\psi \cup \{(e, o)\})) \end{aligned}$$

Similarly, we have

$$f_{wc}(e | \psi') = \min_{o \in O(e, \psi')} p_H(\mathcal{H}(\psi')) - p_H(\mathcal{H}(\psi' \cup \{(e, o)\}))$$

We next show that  $f_{wc}(e | \psi) \geq f_{wc}(e | \psi')$ . We first consider the case when there exists a state  $o' \in \arg \min_{o \in O(e, \psi)} p_H(\mathcal{H}(\psi)) - p_H(\mathcal{H}(\psi \cup \{(e, o)\}))$  such that  $o' \in O(e, \psi')$ . In this case, we have

$$f_{wc}(e | \psi) = p_H(\mathcal{H}(\psi)) - p_H(\mathcal{H}(\psi \cup \{(e, o')\})) \quad (37)$$

and

$$\begin{aligned} f_{wc}(e | \psi') &= \min_{o \in O(e, \psi')} p_H(\mathcal{H}(\psi')) - p_H(\mathcal{H}(\psi' \cup \{(e, o)\})) \\ &\leq p_H(\mathcal{H}(\psi')) - p_H(\mathcal{H}(\psi' \cup \{(e, o')\})) \end{aligned} \quad (38)$$

Because  $\psi \subseteq \psi'$ , any hypotheses  $h \in \mathcal{H}(\psi')$  that is not consistent with  $\psi' \cup \{(e, o')\}$  must satisfy that  $h \in \mathcal{H}(\psi)$  and  $h$  is not consistent with  $\psi \cup \{(e, o')\}$ . Thus, we have  $\mathcal{H}(\psi') \setminus \mathcal{H}(\psi' \cup \{(e, o')\}) \subseteq \mathcal{H}(\psi) \setminus \mathcal{H}(\psi \cup \{(e, o')\})$ . Thus,  $p_H(\mathcal{H}(\psi)) - p_H(\mathcal{H}(\psi \cup \{(e, o')\})) \leq p_H(\mathcal{H}(\psi')) - p_H(\mathcal{H}(\psi' \cup \{(e, o')\}))$ . This together with (37) and (38) implies that  $f_{wc}(e | \psi) \geq f_{wc}(e | \psi')$ .

We next consider the case when there does not exist a state  $o' \in \arg \min_{o \in O(e, \psi)} p_H(\mathcal{H}(\psi)) - p_H(\mathcal{H}(\psi \cup \{(e, o)\}))$  such that  $o' \in O(e, \psi')$ . In this case, for any state  $o'' \in \arg \min_{o \in O(e, \psi)} p_H(\mathcal{H}(\psi)) - p_H(\mathcal{H}(\psi \cup \{(e, o)\}))$ , we have

$$\mathcal{H}(\psi') \subseteq \mathcal{H}(\psi) \setminus \mathcal{H}(\psi \cup \{(e, o'')\}) \quad (39)$$

This is because none of the hypotheses from  $\mathcal{H}(\psi')$  is consistent with  $\{(e, o'')\}$ . It follows that

$$\begin{aligned} f_{wc}(e | \psi) &= p_H(\mathcal{H}(\psi)) - p_H(\mathcal{H}(\psi \cup \{(e, o'')\})) \\ &\geq p_H(\mathcal{H}(\psi')) \\ &\geq \min_{o \in O(e, \psi')} p_H(\mathcal{H}(\psi')) - p_H(\mathcal{H}(\psi' \cup \{(e, o)\})) \\ &= f_{wc}(e | \psi') \end{aligned}$$

The first equality is due to the definition of  $o''$  and the first inequality is due to (39). This finishes the proof of this proposition.  $\square$

### A.3. Proof of Proposition 2

*Proof:* The proof of worst-case monotonicity is trivial. For any  $e \in E$  and  $\psi$  such that  $e \notin \text{dom}(\psi)$ , let  $o' \in \arg \min_{o \in O(e, \psi)} f(\text{dom}(\psi) \cup \{e\}, \psi \cup \{(e, o)\}) - f(\text{dom}(\psi))$ . Consider any realization  $\phi'$  that is consistent with  $\psi \cup \{(e, o')\}$ , i.e.,  $\phi' \sim \psi \cup \{(e, o')\}$ , we have

$$\begin{aligned} f_{wc}(e | \psi) &= f(\text{dom}(\psi) \cup \{e\}, \psi \cup \{(e, o')\}) - f(\text{dom}(\psi), \psi) \\ &= f(\text{dom}(\psi) \cup \{e\}, \phi') - f(\text{dom}(\psi), \phi') \\ &\geq 0 \end{aligned}$$

The second equality is due to  $f : 2^E \times O^E \rightarrow \mathbb{R}_{\geq 0}$  satisfies the property of minimal dependency and the inequality is due to  $f : 2^E \times O^E \rightarrow \mathbb{R}_{\geq 0}$  is pointwise monotone with respect to  $p(\phi)$ .

We next prove that  $f : 2^E \times O^E \rightarrow \mathbb{R}_{\geq 0}$  is worst-case submodular with respect to  $p(\phi)$ . Consider any two partial realizations  $\psi$  and  $\psi'$  such that  $\psi \subseteq \psi'$  and any item  $e \in E \setminus \text{dom}(\psi')$ . Because of the first condition, we have  $O(e, \psi) = O(e, \psi')$ . Then for any  $o \in O(e, \psi)$ ,  $\phi \sim \psi \cup \{(e, o)\}$  and  $\phi' \sim \psi' \cup \{(e, o)\}$ , we have

$$\begin{aligned} &f(\text{dom}(\psi) \cup \{e\}, \psi \cup \{(e, o)\}) - f(\text{dom}(\psi), \psi) \\ &= f(\text{dom}(\psi) \cup \{e\}, \phi) - f(\text{dom}(\psi), \phi) \\ &= f(\text{dom}(\psi) \cup \{e\}, \phi') - f(\text{dom}(\psi), \phi') \\ &\geq f(\text{dom}(\psi') \cup \{e\}, \phi') - f(\text{dom}(\psi'), \phi') \\ &= f(\text{dom}(\psi') \cup \{e\}, \psi' \cup \{(e, o)\}) - f(\text{dom}(\psi'), \psi') \end{aligned} \tag{40}$$

All equalities are due to  $f : 2^E \times O^E \rightarrow \mathbb{R}_{\geq 0}$  satisfies the property of minimal dependency and the inequality is due to  $f : 2^E \times O^E \rightarrow \mathbb{R}_{\geq 0}$  is pointwise submodular.

Let  $o' \in \arg \min_{o \in O(e, \psi)} f(\text{dom}(\psi) \cup \{e\}, \psi \cup \{(e, o)\}) - f(\text{dom}(\psi), \psi)$ , we have

$$\begin{aligned} f_{wc}(e | \psi) &= f(\text{dom}(\psi) \cup \{e\}, \psi \cup \{(e, o')\}) - f(\text{dom}(\psi), \psi) \\ &\geq f(\text{dom}(\psi') \cup \{e\}, \psi \cup \{(e, o')\}) - f(\text{dom}(\psi'), \psi') \\ &\geq \min_{o \in O(e, \psi')} f(\text{dom}(\psi') \cup \{e\}, \psi \cup \{(e, o)\}) - f(\text{dom}(\psi'), \psi') \\ &= f_{wc}(e | \psi') \end{aligned} \tag{41}$$

The first inequality is due to (40). This finishes the proof of this proposition.  $\square$

### A.4. Proof of Proposition 3

*Proof:* It is easy to verify that  $f : 2^E \times O^E \rightarrow \mathbb{R}_{\geq 0}$  is worst-case monotone and it satisfies the property of minimal dependency. We next show that it is also worst-case submodular with respect to  $p(\phi)$ .

Consider any two partial realizations  $\psi$  and  $\psi'$  such that  $\psi \subseteq \psi'$  and any node  $e \in E \setminus \text{dom}(\psi')$ . Let  $Z(\psi)$  denote the set of edges whose status is observed under  $\psi$  and let  $E(\psi)$  denote the set of nodes that are influenced under  $\psi$ . Clearly,  $Z(\psi) \subseteq Z(\psi')$  and  $E(\psi) \subseteq E(\psi')$ . Now consider the worst-case marginal utility  $f_{wc}(e | \psi)$  of  $e$  on top of  $\psi$ , the worst-case realization of  $e$ , say  $o$ , occurs when all edges from  $\{(u, v) \in Z \setminus Z(\psi) \mid p(u, v) < 1\}$  are blocked. Similarly, for the worst-case marginal utility  $f_{wc}(e | \psi')$  of  $e$  on top of  $\psi'$ ,

the worst-case realization of  $e$ , say  $o'$ , occurs when all edges from  $\{(u, v) \in Z \setminus Z(\psi') \mid p(u, v) < 1\}$  are blocked. Because  $Z(\psi) \subseteq Z(\psi')$ , we have  $Z \setminus Z(\psi') \subseteq Z \setminus Z(\psi)$ . This together with  $E(\psi) \subseteq E(\psi')$  implies that if a node  $w \in E \setminus E(\psi')$  can be reached by some live edge under  $(e, o')$ , it must be the case that  $w \in E \setminus E(\psi)$  and  $w$  can be reached by some live edge under  $(e, o)$ . We conclude that the additional nodes influenced by  $e$  on top of  $\psi$  under the worst-case realization is a superset of the additional nodes influenced by  $e$  on top of  $\psi'$  under the worst-case realization. It follows that  $f_{wc}(e \mid \psi) \geq f_{wc}(e \mid \psi')$ . This finishes the proof of this proposition.  $\square$