

## Notation and proofs

In this e-companion we present the notation used throughout this paper, and we complete the proofs of some propositions.

### EC.1. Notation

#### Sets

- $I$ : set of demand nodes;
- $J$ : set of candidate facility locations (leader and competition);
- $J_c$ : set of competition's facilities;
- $J_1$ : set of leader's candidate sites;
- $J_1^* \subseteq J_1$ : set of leader's open facilities
- $J^* \subseteq J$ : set of open facilities (leader and competitor).

#### Parameters

- $d_i$ : demand originating from node  $i \in I$ ;
- $t_{ij}$ : travel time between nodes  $i \in I$  and  $j \in J$ ;
- $\alpha$ : coefficient of the waiting time in the disutility formula;
- $\beta$ : coefficient of the balking probability in the disutility formula;
- $B$ : available budget (for opening new facilities and associated service rates);
- $c_f$ : fixed cost associated with opening a new facility;
- $c_\mu$ : cost per unit of service;
- $\bar{\mu}$ : maximum service rate allowed by the budget;
- $p$ : number of facilities to open.

#### Basic decision variables

- $y_j$ : binary variable set to 1 if a facility is open at site  $j$ , and to 0 otherwise;
- $\mu_j$ : service rate at open facilities.

### Additional variables

$x_{ij}$ : arrival rate at facility  $j \in J$  originating from demand node  $i \in I$ ;

$\lambda_j$ : arrival rate at node  $j \in J$ ;

$\rho_j$ : utilization rate of facility  $j \in J$ ;

$\bar{\lambda}_j$ : throughput rate (customers accessing service) at node  $j \in J$ ;

$w_j$ : mean queueing time at facility  $j$ .

## EC.2. Proofs of Propositions 1, 2, 4, 5, 6, 7, 9, 10, 11 and Theorem 1

PROPOSITION 1. *The waiting time  $w_j$  is increasing in  $\lambda_j$ .*

*Proof.* The derivative of  $w_j$  with respect to  $\lambda_j$  (see Equation (10)) is

$$\frac{\partial w_j}{\partial \lambda_j} = \frac{\partial w_j}{\partial \rho_j} \frac{\partial \rho_j}{\partial \lambda_j} = \frac{\partial w_j}{\partial \rho_j} \frac{1}{\mu_j}.$$

To show that  $\partial w_j / \partial \rho_j$  is nonnegative for all  $\rho_j \neq 1$ , let us consider

$$\frac{\partial w_j}{\partial \rho_j} = \frac{1}{\mu_j} \left( -\frac{K^2 \rho_j^{K-1}}{(\rho_j^K - 1)^2} + \frac{1}{(\rho_j - 1)^2} \right).$$

Basic algebraic manipulation yields

$$\frac{1}{(\rho_j - 1)^2} \geq \frac{K^2 \rho_j^{K-1}}{(\rho_j^K - 1)^2} \iff \sum_{i=0}^{K-1} \rho_j^i \geq K \rho_j^{(K-1)/2}. \quad (\text{EC.1})$$

To prove that the right-hand inequality holds true, we consider two cases.

If  $K$  is odd:

$$\begin{aligned} \sum_{i=0}^{K-1} \rho_j^i &= \sum_{i=0}^{(K-1)/2-1} (\rho_j^i + \rho_j^{K-1-i}) + \rho_j^{(K-1)/2} \\ &\geq 2 \sum_{i=0}^{(K-1)/2-1} \rho_j^{(K-1)/2} + \rho_j^{(K-1)/2} = K \rho_j^{(K-1)/2}. \end{aligned}$$

If  $K$  is even:

$$\sum_{i=0}^{K-1} \rho_j^i = \sum_{i=0}^{(K-2)/2} \left( \rho_j^i + \rho_j^{K-1-i} \right) \geq 2 \sum_{i=0}^{(K-2)/2} \rho_j^{(K-1)/2} = K \rho_j^{(K-1)/2}.$$

It follows that  $w_j$  is an increasing function of  $\lambda_j$ .

**PROPOSITION 2.** *The probability of balking  $p_{Kj}$  is increasing in  $\lambda_j$ .*

*Proof.* The derivative of  $p_{Kj}$  with respect to  $\lambda_j$  is

$$\begin{aligned} p'_{Kj} &= \frac{\lambda_j^{K-1} \mu_j}{\left( \lambda_j^{K+1} - \mu_j^{K+1} \right)^2} \left[ \lambda_j^{K+1} - (K+1) \lambda_j \mu_j^K + K \mu_j^{K+1} \right] \\ &= \sigma [x^{K+1} - (K+1)x + K], \end{aligned}$$

where  $\sigma$  is a positive number and  $x = \lambda_j / \mu_j$ . By differentiating with respect to  $x$ , we find that the right-hand-side achieves its minimum value 0 at  $x = 1$ , which concludes the proof.

**PROPOSITION 4.** *When  $K = \infty$ , i.e., balking does not occur (in this case, the model admits a solution only if the total service rate exceeds the total demand rate), the lower level objective function is convex jointly in  $\lambda$  and  $\mu$ .*

*Proof.* If  $K = \infty$ , the probability of balking can be removed from the objective, since it is equal to 0. Moreover,  $w_j = 1/(\mu_j - \lambda_j)$ , and the lower level objective takes the form

$$\sum_{i \in I} \sum_{j \in J^*} \left[ \frac{1}{\theta} x_{ij} \ln x_{ij} + x_{ij} t_{ij} \right] - \alpha \sum_{j \in J^*} \ln(\mu_j - \lambda_j).$$

Basic algebra shows that its Hessian is positive semidefinite, hence the function is convex.

**PROPOSITION 5.** *The integral of the waiting time,  $W_j(\lambda_j, \mu_j)$  is pseudoconvex.*

*Proof.* Let  $x = (\lambda_x, \mu_x)$  and  $y = (\lambda_y, \mu_y)$ . Assume that  $\nabla W(x)(y - x) \geq 0$ . Then we have:

$$\begin{aligned} &\left( w_j(x), -\frac{\lambda_x}{\mu_x} w_j(x) \right) (\lambda_y - \lambda_x, \mu_y - \mu_x) \geq 0 \\ \Rightarrow &(\rho_y - \rho_x) w_j(x) \geq 0 \\ \Rightarrow &\rho_y \geq \rho_x, \end{aligned} \tag{EC.2}$$

since  $w_j$  is nonnegative. On the other hand,  $\partial W_j/\partial \rho = \mu_j w_j$  is nonnegative, we have that  $W_j$  is increasing in  $\rho$ , so  $\rho_y \geq \rho_x \Rightarrow W_j(y) \geq W_j(x)$ . From Eq (EC.2) it follows that if  $\nabla W(x)(y-x) \geq 0$  then  $W_j(y) \geq W_j(x)$ , hence  $W_j$  is pseudoconvex.

PROPOSITION 6.  $G_1$  is strongly monotone in  $x$  of modulus  $\theta \cdot d_{MAX}$ .

*Proof.* Gilbert et al. (2015) have already argued that  $G_1$  is strongly monotone. Indeed, the associated Jacobian is a positive definite diagonal matrix over  $D$ , with the smallest possible eigenvalue  $1/(\theta \cdot d_{MAX})$ . It follows that  $G_1$  is strongly monotone with modulus  $\theta \cdot d_{MAX}$ .

PROPOSITION 7.  $G_2$  is monotone in  $x$ .

*Proof.*

$$\begin{aligned}
\langle G_2(\mu, x) - G_2(\mu, y), x - y \rangle &= \sum_{i \in I} \sum_{j \in J^*} \left( \frac{1}{\mu_j - \sum_{l \in I} x_{l,j}} - \frac{1}{\mu_j - \sum_{l \in I} y_{l,j}} \right) \cdot (x_{ij} - y_{ij}) \\
&= \sum_{j \in J^*} \left[ \frac{\mu_j - \sum_{l \in I} y_{l,j} - \mu_j + \sum_{l \in I} x_{l,j}}{\left( \mu_j - \sum_{l \in I} x_{l,j} \right) \cdot \left( \mu_j - \sum_{l \in I} y_{l,j} \right)} \cdot \sum_{i \in I} (x_{ij} - y_{ij}) \right] \\
&= \sum_{j \in J^*} \left[ \frac{\sum_{l \in I} (x_{l,j} - y_{l,j}) \sum_{l \in I} (x_{l,j} - y_{l,j})}{\left( \mu_j - \sum_{l \in I} x_{l,j} \right) \cdot \left( \mu_j - \sum_{l \in I} y_{l,j} \right)} \right] \\
&\geq 0
\end{aligned}$$

PROPOSITION 9. If  $K = \infty$  and there are no fixed costs, the surrogate model is convex.

*Proof.* According to Proposition 4, the objective is jointly convex in  $\mu$  and  $\lambda$ . Moreover one can, without loss of generality, open all facilities and hence dispense with the binary vector  $y$ . Notwithstanding, a facility can be closed by setting its service level to zero.

PROPOSITION 10. At the optimum of  $(PH^*)$ , if  $K = \infty$ , queueing delays are equal for all leader's facilities.

*Proof.* For fixed  $y$  variables, Equation (66) can be rewritten as

$$\sum_{j \in J^*} \mu_j \leq \bar{\mu}, \quad (\text{EC.3})$$

where  $\bar{\mu}$  is the maximum possible total service rate allowed by the budget. But  $K = \infty$ , so  $w_j(\lambda_j, \mu_j) = 1/(\mu_j - \lambda_j)$  and  $p_{Kj}(\lambda_j, \mu_j) = 0$ , which yields the mathematical program

$$\begin{aligned} (\text{PHY}^*) \quad & \min_{\mu, x} \quad \sum_{i \in I} \sum_{j \in J^*} x_{ij} t_{ij} - \alpha \sum_{j \in J^*} \ln(\mu_j - (\sum_{i \in I} x_{ij})) \\ & \text{s.t.} \quad \text{constraints (65), (68), (69), (EC.3)} \end{aligned}$$

Let  $\delta_i$ ,  $\pi_{ij}$  and  $\gamma$  be the Lagrange multipliers associated with Equations (65), (69) and (EC.3), respectively. Variables  $\delta_i$  are free, while  $\gamma$  and  $\pi_{ij}$  are restricted to be nonnegative. The stationarity conditions of the above program are:

$$\frac{\partial L}{\partial x_{ij}} = 0 \Rightarrow \quad t_{ij} + \alpha w_j(\lambda_j, \mu_j) - \delta_i - \pi_{ij} = 0, \quad \forall i \in I, \forall j \in J^* \quad (\text{EC.4})$$

$$\frac{\partial L}{\partial \mu_j} = 0 \Rightarrow \quad -\alpha w_j(\lambda_j, \mu_j) + \gamma = 0, \quad \forall j \in J^* \cap J_1, \quad (\text{EC.5})$$

and the conclusion follows from Equation (EC.5).

We observe, after plugging  $\alpha w_j(\lambda_j, \mu_j)$  from Equation (EC.5) into Equation (EC.4) for a given demand node  $i$ , that only one flow  $x_{ij}$  is nonzero, provided that transportation times to the leader's facilities are distinct.

**PROPOSITION 11.** *There exists a value of  $\xi^*$  for which  $(\text{PHY}^*(\xi^*))$  yields an optimal solution for  $(P^*)$ .*

*Proof.* Let  $y^*$  and  $\mu^*$  be optimal for  $(P^*)$ . Without loss of generality (there are no fixed costs) we assume that all facilities are open. At equilibrium, let  $c_i^*$  be the cost associated with demand node  $i$  and optimal service rate  $\mu^*$ . Let  $x^*$ ,  $w_j(x^*, \mu_j^*)$  and  $c_i^*$  satisfy Equation (24) and (25). If  $x_{ij}$  is positive, we have:

$$t_{ij} + \alpha w_j(x^*, \mu^*) = c_i^*, \quad \forall j \in J, \forall i \in I. \quad (\text{EC.6})$$

Let  $C = \max_{i \in I} \{c_i^*\}$  in the initial formulation. For  $j \in J$ , we let  $\xi_j = c_i^* - t_{ij} - C$  and select an index  $i$  corresponding to a positive flow  $x_{ij}^*$ . If no such  $i$  exists, then  $\mu_j^* = 0$ , otherwise the leader would waste monetary resources. We then set  $\xi_j = -C$ .

Now, let  $\delta_i$ ,  $\pi_{ij}$  and  $\gamma$  be the Lagrange multipliers associated with Equations (65), (69) and (EC.3), respectively. Variables  $\delta_i$  and  $\gamma$  are free, while  $\pi_{ij}$  are restricted to be nonnegative. The stationarity conditions of the program above take the form

$$\frac{\partial \mathcal{L}}{\partial x_{ij}} = 0 \Rightarrow \quad t_{ij} + \alpha w_j(x, \mu_j) - \delta_i = 0, \quad \text{if } x_{ij} > 0, \quad \forall i \in I, \forall j \in J \quad (\text{EC.7})$$

$$\frac{\partial \mathcal{L}}{\partial \mu_j} = 0 \Rightarrow \quad -\alpha w_j(x, \mu_j) + \gamma + \xi_j = 0, \quad \forall j \in J_1. \quad (\text{EC.8})$$

Note that the derivative of the Lagrangian with respect to  $x_{ij}$  is left unchanged, i.e., Equation (EC.7) is equivalent to Equation (EC.4). If  $\gamma = C$ , we derive from Equation (EC.8) that  $\alpha w_j(x, \mu_j) = c_i^* - t_{ij}$ , which is equivalent to Equation (EC.7). This completes the proof, since for the given values of  $\xi$ , variables  $x$  and  $\mu$  match the optimal solution of (P\*).

**THEOREM 1.** *The error of the upper-level objective function is  $O(1/N_1 + 1/N_2)$ , where  $N_1$  and  $N_2$  are the number of samples for the linearization of  $g_1$  and  $g_2$ , respectively.*

*Proof.* Let  $\bar{G}$  be an approximation of  $G$ . We denote by  $\bar{x}$  the solution of  $\text{IV}(\bar{G}(\mu, \cdot), D)$ , and by  $x$  the solution of  $\text{IV}(G(\mu, \cdot), D)$ . Then the following inequalities hold:

$$\begin{aligned} \langle G(\mu, x), \bar{x} - x \rangle &\geq 0 \\ \langle \bar{G}(\mu, \bar{x}), x - \bar{x} \rangle &\geq 0 \\ \Rightarrow \langle \bar{G}(\mu, \bar{x}) - G(\mu, x), x - \bar{x} \rangle &\geq 0 \end{aligned} \quad (\text{EC.9})$$

From the strong monotonicity of  $G$  and Eq. (EC.9) it follows that

$$\langle \bar{G}(\mu, \bar{x}) - G(\mu, \bar{x}), x - \bar{x} \rangle \geq \frac{1}{\theta \cdot d_{\text{MAX}}} \|x - \bar{x}\|^2. \quad (\text{EC.10})$$

Applying the Cauchy-Schwarz inequality, we obtain

$$\theta \cdot d_{\text{MAX}} \cdot \|\bar{G}(\mu, \bar{x}) - G(\mu, \bar{x})\| \geq \|x - \bar{x}\|. \quad (\text{EC.11})$$

It follows that

$$\begin{aligned} |f(x) - f(\bar{x})| &= \left| \sum_{i \in I} \sum_{j \in J_1^*} (x_{ij} - \bar{x}_{ij}) \right| \leq \sqrt{|I| \cdot |J|} \|x - \bar{x}\| \quad (\text{Cauchy-Schwarz inequality}) \\ &\leq \sqrt{|I| \cdot |J|} \theta \cdot d_{\text{MAX}} \cdot \|\bar{G}(\mu, \bar{x}) - G(\mu, \bar{x})\|. \end{aligned} \quad (\text{EC.12})$$

We perform two separate linear approximations on  $g_1$  and  $g_2$ , respectively. Then the mappings  $\bar{G}_1(x)$  and  $\bar{G}_2(x)$  are piecewise constant approximations, that we detail separately.

A.  $\bar{G}_1$ : Each component  $(i, j)$  of this vector is a piecewise constant approximation of  $\log(x_{ij})$ , satisfying:

- i) there are  $N_1$  total samples on  $x_{ij}$ , starting from  $r_{\min}$  to  $d_{\text{MAX}}$ ;
- ii) the sampling points are chosen so that the segments are vertically equidistant;
- iii) the vertical positions of the segments are the slopes of the tangents to  $x \log(x)$ , evaluated at the sampling points.

Let  $\Delta_1$  be the difference between two consecutive slope values:

$$\Delta_1 = \frac{\log(d_{\text{MAX}}) - \log(r_{\min})}{N_1 - 1}.$$

Then  $|\bar{G}_{1(i,j)} - G_{1(i,j)}| \leq \Delta_1$ , which yields:

$$\|\bar{G}_1(x) - G_1(x)\| = \sqrt{\sum_{i \in I} \sum_{j \in J^*} |\bar{G}_{1(i,j)} - G_{1(i,j)}|^2} \leq \frac{(\log(d_{\text{MAX}}) - \log(r_{\min})) \sqrt{|I| \cdot |J|}}{N_1 - 1} \quad (\text{EC.13})$$

B.  $\bar{G}_2$ : is a mapping whose  $(i, j)$ -component is a constant piecewise approximation of  $1/q_j$ , where  $q_j = \mu_j - \sum_{i \in I} x_{ij}$ . Similar to  $\bar{G}_1$ , this linearization satisfies the following:

- i) there are  $N_2$  total samples, starting from  $\psi$  to  $\mu_{\text{MAX}}$ ;
- ii) the sampling points are chosen so that the segments are vertically equidistant;
- iii) the vertical positions of the segments are the slopes of the tangents  $-\log(q)$  evaluated at the sampling points.

We note by  $\Delta_2$  the difference between two consecutive slope values:

$$\Delta_2 = \frac{\frac{1}{\psi} - \frac{1}{\mu_{\text{MAX}}}}{N_2 - 1} /$$

Then  $|\bar{G}_{2(i,j)} - G_{2(i,j)}| \leq \Delta_2$ , which yields

$$\|\bar{G}_2(x) - G_2(x)\| = \sqrt{\sum_{i \in I} \sum_{j \in J^*} |\bar{G}_2(x_{ij}) - G_2(x_{ij})|^2} \leq \frac{\left(\frac{1}{\psi} - \frac{1}{\mu_{\text{MAX}}}\right) \sqrt{|I| \cdot |J|}}{N_2 - 1} \quad (\text{EC.14})$$

From Eq. (EC.12) it follows that, given  $y$  and  $\mu$ :

$$|f(x) - f(\bar{x})| \leq \theta \cdot d_{\text{MAX}} |I| \cdot |J| \left[ \frac{(\log(d_{\text{MAX}}) - \log(r_{\text{min}}))}{N_1 - 1} + \alpha \frac{\frac{1}{\psi} - \frac{1}{\mu_{\text{MAX}}}}{N_2 - 1} \right] \in O\left(\frac{1}{N_1} + \frac{1}{N_2}\right). \quad (\text{EC.15})$$

Theorem 1 has several implications.

- For a given set of open facilities, the absolute difference between the optimal and the approximated objective value is bounded by the right-hand-side of inequality (EC.15). For large values of  $N_1$  and  $N_2$ , the two values are very close.

- If the optimal solution is unique in terms of the location vector  $y$ , and the absolute difference between the objective and other solutions objectives are lower than the right hand side of inequality (EC.15), the approximation algorithm will find the optimum locations.

### EC.3. Linearization of optimality conditions

#### EC.3.1. Complementarity constraints for Program (P2-lin)

Let  $\gamma_i$ ,  $\delta_j$ ,  $\nu_{ij}^n$ ,  $\pi_j^{rp}$ ,  $\eta_j^{rp}$ , and  $\phi_{ij}$  be the dual variables associated with constraints (47), (48), (49), (50), (51) and (52), respectively. Then the complementarity constraints for program (P2-lin) can be written as:

$$\gamma_i \left( \sum_{j \in J^*} x_{ij} - d_i \right) = 0 \quad \forall i \in I \quad (\text{EC.16})$$

$$\delta_j \left( \lambda_j - \sum_{i \in I} x_{ij} \right) = 0 \quad \forall j \in J^* \quad (\text{EC.17})$$

$$\nu_{ij}^n (v_{ij} - a_f^n x_{ij} - b_f^n) = 0 \quad \forall i \in I; \quad \forall j \in J^*; \quad \forall n \in N \quad (\text{EC.18})$$

$$\pi_j^{rp} (u_j - a_g^{rp} \lambda_j - b_g^{rp} \mu_j - c_g^{rp}) = 0 \quad \forall j \in J^*; \quad \forall r \in R; \quad \forall p \in P \quad (\text{EC.19})$$

$$\eta_j^{rp} (z_j - a_h^{rp} \lambda_j - b_h^{rp} \mu_j - c_h^{rp}) = 0 \quad \forall j \in J^*; \quad \forall r \in R; \quad \forall p \in P \quad (\text{EC.20})$$

$$\phi_{ij} x_{ij} = 0 \quad \forall i \in I; \quad \forall j \in J^*, \quad (\text{EC.21})$$

and can be linearized in the standard fashion, through the introduction of binary variables and big-M constants. For instance, the last constraint is replaced by the inequalities

$$\begin{aligned} \phi_{ij} &\leq M u_{ij} \\ x_{ij} &\leq M(1 - u_{ij}), \end{aligned}$$

where  $u_{ij} \in \{0, 1\}$ . Although it is possible to find a valid upper bound for the variable  $\phi_{ij}$ , a large value of  $M$  is required, which leads to a poor relaxation and consequently an ill-behaved branch-and-bound algorithm.

### EC.3.2. Equality between primal and dual objectives

Alternatively, constraints (EC.16) – (EC.21) can be replaced with constraint (EC.22), which represents the equality between the primal and dual objective of (P2-lin). Then the optimality constraints of (P2-lin) are

$$\begin{aligned} &\sum_{i \in I} \gamma_i d_i + \sum_{n \in N} \sum_{i \in I} \sum_{j \in J} v_{ij}^n b_f^n + \sum_{r \in R} \sum_{p \in P} \sum_{j \in J} (b_g^{rp} \mu_j \pi_j^{rp} + b_h^{rp} \mu_j \eta_j^{rp} + c_g^{rp} \pi_j^{rp} + c_h^{rp} \eta_j^{rp}) \\ &= \sum_{i \in I} \sum_{j \in J} \left[ \frac{1}{\theta} v_{ij} + x_{ij} t_{ij} \right] + \alpha \sum_{j \in J} u_j + \beta \sum_{j \in J} z_j, \end{aligned} \quad (\text{EC.22})$$

$$\sum_{j \in J} x_{ij} = d_i, \quad \forall i \in I$$

$$\lambda_j = \sum_{i \in I} x_{ij}, \quad \forall j \in J$$

$$v_{ij} - a_f^n x_{ij} \geq b_f^n, \quad \forall i \in I; \forall j \in J; \forall n \in N$$

$$u_j - a_g^{rp} \lambda_j - b_g^{rp} \mu_j \geq c_g^{rp}, \quad \forall j \in J; \forall r \in R; \forall p \in P$$

$$z_j - a_h^{rp} \lambda_j - b_h^{rp} \mu_j \geq c_h^{rp}, \quad \forall j \in J; \forall r \in R; \forall p \in P$$

$$\begin{aligned}
\gamma_i + \delta_j - \sum_{n \in N} a_f^n \nu_{ij}^n &\leq t_{ij}, & \forall i \in I; \forall j \in J \\
-\delta_j - \sum_{r \in R} \sum_{p \in P} (a_g^{rp} \pi_j^{rp} + a_h^{rp} \eta_j^{rp}) &= 0, & \forall j \in J \\
\sum_{n \in N} \nu_{ij}^n &= \frac{1}{\theta}, & \forall i \in I; \forall j \in J \\
\sum_{r \in R} \sum_{p \in P} \pi_j^{rp} &= \alpha, & \forall j \in J \\
\sum_{r \in R} \sum_{p \in P} \eta_j^{rp} &= \beta, & \forall j \in J \\
\pi_j^{rp}, \eta_j^{rp} &\geq 0, & \forall j \in J; \forall r \in R; \forall p \in P \\
x_{ij} \geq 0, \nu_{ij}^n &\geq 0, & \forall i \in I; \forall j \in J; \forall n \in N.
\end{aligned}$$

To obtain a MILP formulation, we linearize the nonlinear terms  $\mu_j \pi_j^{rp}$  and  $\mu_j \eta_j^{rp}$  via the triangle method described in D'Ambrosio et al. (2010). For each term  $\mu_j \pi_j^{kq}$  we introduce  $2(R-1)(P-1)$  binary variables  $\bar{l}_{jrpq}^\pi$  and  $\underline{l}_{jrpq}^\pi$  associated with the upper and lower triangles, respectively, of the rectangle defined by the intervals  $[\pi^r, \pi^{r+1})$  and  $[\mu^p, \mu^{p+1})$ . Note that the values of  $\pi$  and  $\eta$  are upper bounded by  $\alpha$  and  $\beta$ , respectively. Additionally,  $\mu$  is bounded by the maximum value allowed by the leader's budget,  $\bar{\mu}$ . Next, we introduce  $J_1 R P$  continuous variables  $s_{jrpq} \in [0, 1]$  which will be used to express the couple  $(\pi_j^{kq}, \mu_j)$  as a convex combination of triangle vertices. We introduce a similar linearization for the term  $\mu_j \eta_j^{kq}$ . The approximation of  $\mu_j \pi_j^{kq}$  and  $\mu_j \eta_j^{kq}$  is then

$$\sum_{r=1}^{R-1} \sum_{p=1}^{P-1} (\bar{l}_{jrpq}^\pi + \underline{l}_{jrpq}^\pi) = 1, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.23})$$

$$\begin{aligned}
s_{jrpq}^\pi &\leq \bar{l}_{jrpq}^\pi + \underline{l}_{jrpq}^\pi + \bar{l}_{jrp-1kq}^\pi + \underline{l}_{jr-1p-1kq}^\pi + \bar{l}_{jr-1p-1kq}^\pi + \underline{l}_{jr-1pkq}^\pi, \\
&\forall j \in J_1; \forall r \in R; \forall p \in P; \forall k \in R; \forall q \in P \quad (\text{EC.24})
\end{aligned}$$

$$\sum_{r=1}^R \sum_{p=1}^P s_{jrpq}^\pi = 1, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.25})$$

$$\pi_j^{kq} = \sum_{r=1}^R \sum_{p=1}^P s_{jrpq}^\pi \pi^r, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.26})$$

$$\mu_j = \sum_{r=1}^R \sum_{p=1}^P s_{jrpq}^\pi \mu^p, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.27})$$

$$e_{jkq}^\pi = \sum_{r=1}^R \sum_{p=1}^P s_{jrpq}^\pi \pi^r \mu^p, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.28})$$

$$\sum_{r=1}^{R-1} \sum_{p=1}^{P-1} \left( \bar{l}_{jrpq}^\eta + \underline{l}_{jrpq}^\eta \right) = 1, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.29})$$

$$s_{jrpq}^\eta \leq \bar{l}_{jrpq}^\eta + \underline{l}_{jrpq}^\eta + \bar{l}_{jrp-1kq}^\eta + \underline{l}_{jrp-1kq}^\eta + \bar{l}_{jr-1p-1kq}^\eta + \underline{l}_{jr-1p-1kq}^\eta, \quad \forall j \in J_1; \forall r \in R; \forall p \in P; \forall k \in R; \forall q \in P \quad (\text{EC.30})$$

$$\sum_{r=1}^R \sum_{p=1}^P s_{jrpq}^\eta = 1, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.31})$$

$$\eta_j^{kq} = \sum_{r=1}^R \sum_{p=1}^P s_{jrpq}^\eta \eta^r, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.32})$$

$$\mu_j = \sum_{r=1}^R \sum_{p=1}^P s_{jrpq}^\eta \mu^p, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.33})$$

$$e_{jkq}^\eta = \sum_{r=1}^R \sum_{p=1}^P s_{jrpq}^\eta \eta^r \mu^p, \quad \forall j \in J_1; \forall k \in R; \forall q \in P \quad (\text{EC.34})$$

The complete MILP formulation is presented below. It involves variables associated with the original fixed point (or bilevel) formulation  $(y, \mu, x)$ , together with variables issued from the linearizations and primal-dual optimality conditions.

$$\begin{aligned}
(\text{P-lin}) \quad & \max_{x, y, \mu, \lambda, u, v, z, e, \pi, \eta, \nu, \gamma, \delta, \bar{l}^\pi, \underline{l}^\pi, s^\pi, e^\pi, s^\eta, \bar{l}^\eta, \underline{l}^\eta, s, \bar{l}, \underline{l}} \sum_{j \in J_1} e_j \\
& \sum_{j \in J_1} y_j c_f + \sum_{j \in J_1} c_\mu \mu_j \leq B, \\
& \mu_j \leq \bar{\mu} y_j, \quad \forall j \in J_1 \\
& \sum_{r \in R} \sum_{p \in P} \sum_{j \in J_c} (b_g^{rp} \pi_j^{rp} \mu_j + b_h^{rp} \eta_j^{rp} \mu_j + c_g^{rp} \pi_j^{rp} + c_h^{rp} \eta_j^{rp}) + \sum_{n \in N} \sum_{i \in I} \sum_{j \in J} \nu_{ij}^n b_f^n \\
& + \sum_{r \in R} \sum_{p \in P} \sum_{j \in J_1} (b_g^{rp} e_{jrp}^\pi + b_h^{rp} e_{jrp}^\eta + c_g^{mp} \pi_j^{rp} + c_h^{rp} \eta_j^{rp}) + \sum_{i \in I} \gamma_i d_i \\
& = \sum_{i \in I} \sum_{j \in J} \left[ \frac{1}{\theta} v_{ij} + x_{ij} t_{ij} \right] + \alpha \sum_{j \in J} u_j + \beta \sum_{j \in J} z_j, \\
& \sum_{j \in J} x_{ij} = d_i, \quad \forall i \in I
\end{aligned}$$

$$\begin{aligned}
\lambda_j &= \sum_{i \in I} x_{ij}, & \forall j \in J \\
v_{ij} - a_f^n x_{ij} &\geq b_f^n, & \forall i \in I; \forall j \in J; \forall n \in N \\
u_j - a_g^{rp} \lambda_j - b_g^{rp} \mu_j &\geq c_g^{rp}, & \forall j \in J; \forall r \in R; \forall p \in P \\
z_j - a_h^{rp} \lambda_j - b_h^{rp} \mu_j &\geq c_h^{rp}, & \forall j \in J; \forall r \in R; \forall p \in P \\
\gamma_i + \delta_j - \sum_{n \in N} a_f^n v_{ij}^n &\leq t_{ij}, & \forall i \in I; \forall j \in J \\
-\delta_j - \sum_{r \in R} \sum_{p \in P} (a_g^{rp} \pi_j^{rp} + a_h^{rp} \eta_j^{rp}) &= 0, & \forall j \in J \\
\sum_{n \in N} v_{ij}^n &= \frac{1}{\theta}, & \forall i \in I; \forall j \in J \\
\sum_{r \in R} \sum_{p \in P} \pi_j^{rp} &= \alpha, & \forall j \in J \\
\sum_{r \in R} \sum_{p \in P} \eta_j^{rp} &= \beta, & \forall j \in J \\
\text{constraints (EC.23) - (EC.34) and (53) - (58),} & & \\
y_j \in \{0, 1\}, \mu_j, \pi_j^{rp}, \eta_j^{rp} &\geq 0, & \forall j \in J; \forall r \in R; \forall p \in P \\
x_{ij} \geq 0, v_{ij}^n &\geq 0, & \forall i \in I; \forall j \in J; \forall n \in N.
\end{aligned}$$

### EC.3.3. Example of lower level linearization when $K = \infty$

Recall that, according to Proposition 4, the function is convex if the buffer zone is infinite (no balking). In that situation, the maximum of the linear approximations is consistent with the original function, give or take the approximation error. Proceeding as before, we obtain

$$g^{rp}(\lambda, \mu) = a_g^{rp} \lambda + b_g^{rp} \mu + c_g^{rp} = \frac{\alpha}{\mu^p - \lambda^r} \lambda - \frac{\alpha}{\mu^p - \lambda^r} \mu - \alpha(\ln(\mu^p - \lambda^r) - 1). \quad (\text{EC.35})$$

This yields the linearized lower level program

$$(\text{P2}^\infty) \quad \min_{x, v, u, \lambda} \quad \sum_{i \in I} \sum_{j \in J^*} \left[ \frac{1}{\theta} v_{ij} + x_{ij} t_{ij} \right] + \alpha \sum_{j \in J^*} u_j \quad (\text{EC.36})$$

$$\text{s.t.} \quad \sum_{j \in J^*} x_{ij} = d_i, \quad \forall i \in I \quad (\text{EC.37})$$

$$\lambda_j = \sum_{i \in I} x_{ij}, \quad \forall j \in J^* \quad (\text{EC.38})$$

$$v_{ij} - a_f^n x_{ij} \geq b_f^n, \quad \forall i \in I; \forall j \in J^*; \forall n \in N \quad (\text{EC.39})$$

$$u_j - a_g^{rp} \lambda_j - b_g^{rp} \mu_j \geq c_g^{rp}, \quad \forall j \in J^*; \forall r \in R; \forall p \in P \quad (\text{EC.40})$$

$$x_{ij} \geq 0, \quad \forall i \in I; \forall j \in J^*. \quad (\text{EC.41})$$

### EC.3.4. Taxonomy

This section provides a taxonomy of the models most relevant to our research, with respect to four features: (i) user choice environment (yes or no), (ii) stochastic (or not), (iii) inclusion of congestion (or not) at facilities, (iv) inclusion (or not) of competition. The relevant information is displayed in Table EC.1.

Authors	user choice	stochastic	congestion	competition
Abouee-Mehrizi et al. (2011)	×	×	×	
Averbakh et al. (2007)	×			
Berman and Drezner (2006)	×		×	
Camacho-Vallejo et al. (2014)	×			
Castillo et al. (2009)		×	×	
Desrochers et al. (1995)			×	
Kim (2013)			×	
Küçükaydin et al. (2011)	×	×		×
Labbé and Hakimi (1991)				×
Marianov and Serra (2001)			×	
Marianov (2003)		×	×	
Marianov et al. (2008)	×	×	×	×
Marić et al. (2012)	×			
Rahmati et al. (2014)		×	×	
Vidyarthi and Jayaswal (2014)		×	×	
Zhang et al. (2010)	×		×	

**Table EC.1** Taxonomy of congested facility location models