

Online Companion for

“Production and Inventory Models Using  
Net Present Value”

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# Appendix

## A Proofs of Theorems and Lemmas in section 1

### Proof of Lemma 1.2 (Properties of $TCPV(\theta|k, v, w)$ ):

Part 1 can be proved by computing the second derivative and using the inequality  $e^\theta > 1 + \theta$  for all  $\theta > 0$ . Part 2 is an immediate result of Part 1 and the definition of Lambert's W Function:

$$1 + \theta^* + q = e^{\theta^*},$$

$$1 = e^{-\theta^*} (\theta^* + 1 + q),$$

$$-e^{-1-q} = e^{-\theta^*-1-q} (-\theta^* - 1 - q),$$

$$-\theta^* - 1 - q = \text{LambertW}(-1, -e^{-1-q}),$$

$$\theta^* = -1 - q - \text{LambertW}(-1, -e^{-1-q}).$$

Part 3 is easy to verify. □

### Proof of Lemma 1.3

We know that

$$TCPV(\theta^*|k, v, w) = \frac{2}{\alpha^2} (k(\alpha^2/2)f_1(\theta^*) + vf_2(\theta^*) - w) \geq 0,$$

$$TCPV(r\theta^*|k, v, w) = \frac{2}{\alpha^2} (k(\alpha^2/2)f_1(r\theta^*) + vf_2(r\theta^*) - w) \geq 0,$$

and  $TCPV(r\theta^*|k, v, w) \geq TCPV(\theta^*|k, v, w)$  means

$$k(\alpha^2/2)f_1(r\theta^*) + vf_2(r\theta^*) \geq k(\alpha^2/2)f_1(\theta^*) + vf_2(\theta^*). \quad (\text{A.1})$$

It is easy to see that

$$\text{If } a \geq b \geq v, w \geq v, \quad \text{then } \frac{a-w}{b-w} \leq \frac{a-v}{b-v}.$$

Based on inequalities (A.1) and (1.3b), we have the inequality in the following derivation.

$$\begin{aligned}
 \frac{TCPV(r\theta^*|k, v, w)}{TCPV(\theta^*|k, v, w)} &= \frac{k(\alpha^2/2)f_1(r\theta^*) + vf_2(r\theta^*) - w}{k(\alpha^2/2)f_1(\theta^*) + vf_2(\theta^*) - w} \\
 &\leq \frac{k(\alpha^2/2)f_1(r\theta^*) + vf_2(r\theta^*) - v}{k(\alpha^2/2)f_1(\theta^*) + vf_2(\theta^*) - v} = \frac{\frac{k(\alpha^2/2) + vr\theta^*}{1 - e^{-r\theta^*}} - v}{\frac{k(\alpha^2/2) + v\theta^*}{1 - e^{-\theta^*}} - v} \\
 &= \frac{\frac{(k/v)(\alpha^2/2) + r\theta^*}{1 - e^{-r\theta^*}} - 1}{\frac{(k/v)(\alpha^2/2) + \theta^*}{1 - e^{-\theta^*}} - 1} = \frac{((k/v)(\alpha^2/2) + r\theta^* - 1 + e^{-r\theta^*})(1 - e^{-\theta^*})}{((k/v)(\alpha^2/2) + \theta^* - 1 + e^{-\theta^*})(1 - e^{-r\theta^*})}.
 \end{aligned}$$

From equation (1.7), we may replace  $(k/v)(\alpha^2/2)$  by  $e^{\theta^*} - 1 - \theta^*$  and have

$$\begin{aligned}
 \frac{TCPV(r\theta^*|k, v, w)}{TCPV(\theta^*|k, v, w)} &\leq \frac{(e^{\theta^*} - 1 - \theta^* + r\theta^* - 1 + e^{-r\theta^*})(1 - e^{-\theta^*})}{(e^{\theta^*} - 1 - \theta^* + \theta^* - 1 + e^{-\theta^*})(1 - e^{-r\theta^*})} \\
 &= \frac{(e^{-\theta^*} - 2 - \theta^* + r\theta^* + e^{-r\theta^*})(1 - e^{-\theta^*})}{(e^{\theta^*} - 2 + e^{-\theta^*})(1 - e^{-r\theta^*})} \\
 &= \frac{(e^{\theta^*} - 2 - \theta^* + r\theta^* + e^{-r\theta^*})(1 - e^{-\theta^*})}{e^{\theta^*}(1 - e^{-\theta^*})^2(1 - e^{-r\theta^*})} \\
 &= \frac{1 + (r\theta^* - \theta^* - 2 + e^{-r\theta^*})e^{-\theta^*}}{(1 - e^{-\theta^*})(1 - e^{-r\theta^*})}. \tag{A.2}
 \end{aligned}$$

As

$$E(r, \theta^*) = \frac{1 + (r\theta^* - \theta^* - 2 + e^{-r\theta^*})e^{-\theta^*}}{(1 - e^{-\theta^*})(1 - e^{-r\theta^*})}.$$

By using symbolic mathematical tools (such as Maple), we can easily get the following estimate.

$$\frac{E(r, \theta^*)}{\frac{1}{2}\left(r + \frac{1}{r}\right)} \leq (1 + \epsilon), \quad \forall \theta^* > 0 \text{ and } r \in [\sqrt{0.5}, \sqrt{2}], \tag{A.3}$$

where  $\epsilon = 0.000570126$  and the left hand side of (A.3) reaches its maximum value 1.000570126 when  $r = \sqrt{2}$  and  $\theta^* = 0.28768268$ . □

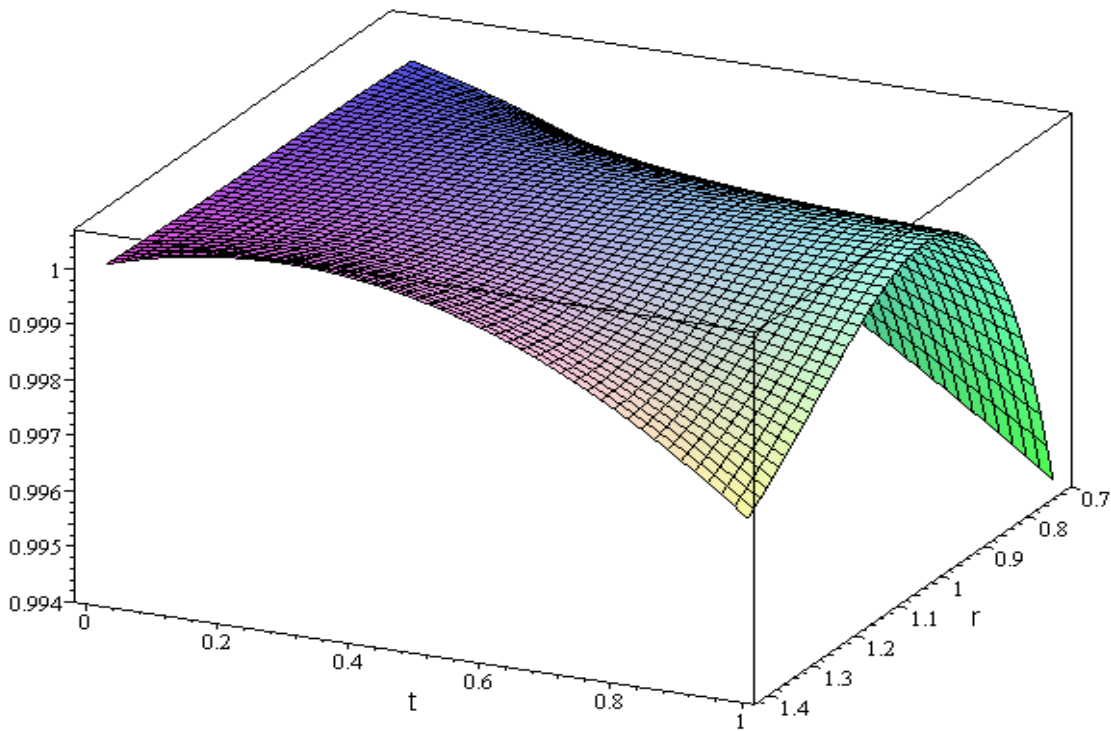
## B Proofs of Theorems and Lemmas in section 3

**Proof of Lemma 3.2 (Minimax Theorem for the continuous relaxation (RP) of (P)):**

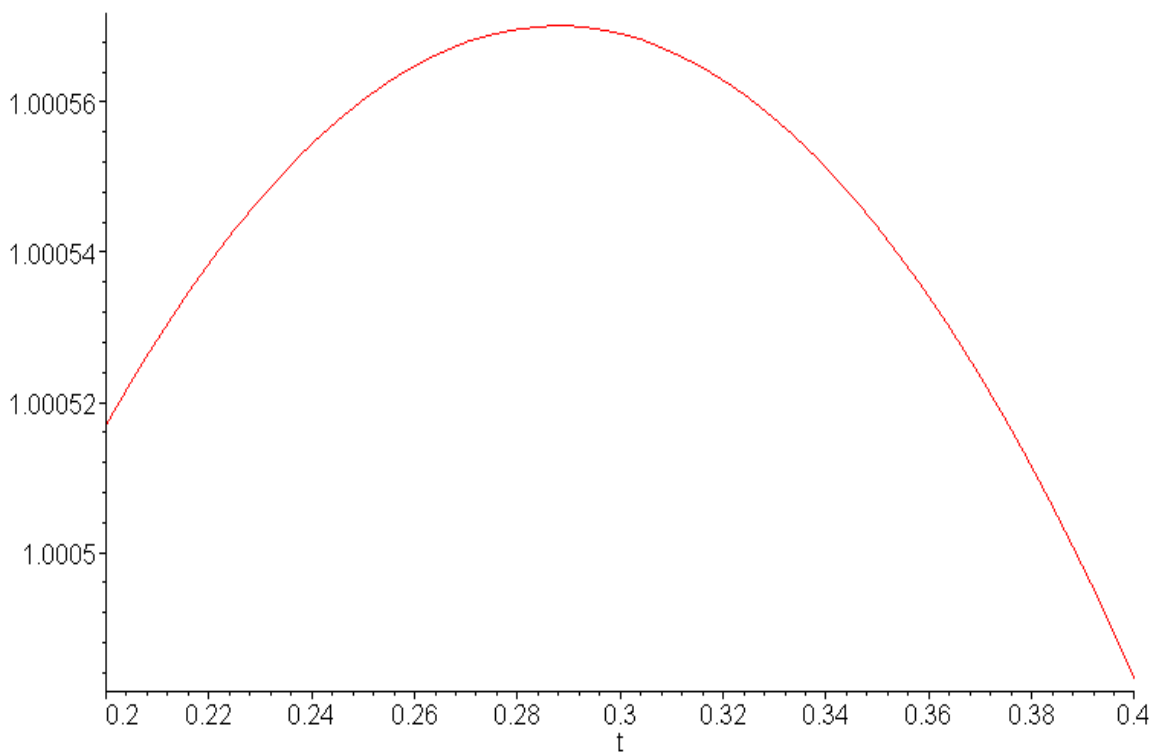
Let

$$\delta \triangleq \min\{K(N) - K(N \setminus \{a\}) : a \in N\}.$$

**Figure 3:** the graph of  $E(r, \theta)/[(r + 1/r)/2]$  for  $r \in [\sqrt{0.5}, \sqrt{2}]$ ,  $\theta > 0$



**Figure 4:** The graph of  $E(r, \theta)/[(r + 1/r)/2]$  for  $r = \sqrt{2}$



Since  $K(\cdot)$  is strictly mono-tone, we have  $\delta > 0$ . Note that for all  $k \in K$  and  $a \in N$ ,  $k_a = K(N) - \sum_{i \neq a} k_i \geq K(N) - K(N \setminus \{a\}) \geq \delta$ . Let  $\bar{T}_L = T_L$  if  $T_L$  is fixed, and  $\bar{T}_L$  be an arbitrary allowable value for  $T_L$  if  $T_L$  is variable. Defining  $t^L \in T$  by  $t_i^L = T_L$  for all  $i \in N$ . we have  $\tau_r^L = \bar{T}_L$  for all  $r \in R$ . Thus for all  $k \in K$  we have

$$f(t^L, k) = K(N)f_1(\alpha\bar{T}_L) + H(N)f_2(\bar{T}_L) = \frac{K(N) + H(N)\bar{T}_L}{1 - e^{-\alpha\bar{T}_L}} \triangleq c^0,$$

where  $H(N) = \sum_{r \in R} \left( u_{i_r} + \frac{h_{i_r}}{\alpha} \right) \frac{d_r}{\alpha}$ . As  $t \searrow 0$ ,  $f_1(t) \rightarrow \infty$ , we may find  $t^-$  satisfy

$$f_1(\alpha t^-) \geq \frac{c^0}{\delta},$$

and  $t^+$  satisfy

$$f_2(\alpha t^+) \triangleq \frac{c^0}{\min_{r \in R} d_r \min_{i \in N_0} \left( \frac{h_i}{\alpha} + u_i \right)}.$$

Define  $\theta = \{t : t^-1 \leq t \leq t^+1, \text{ satisfy (2.4b)}\}$  with  $1 \in R^{|N|}$  a vector of ones. Since  $\theta$  is compact and convex and  $f$  has the properties required in Lemma 1 we have for some  $(t^*, k^*) \in \theta \times K$ :

$$\max_{k \in K} \min_{t \in \theta} TCPV(t, k) = \min_{t \in \theta} \max_{k \in K} TCPV(t, k) = TCPV(t^*, k^*),$$

since

$$\begin{aligned} \frac{c^0}{f_1(\alpha t^-)} \leq \delta &\leq K(N) \leq (K(N) + H(N)\bar{T}_L) = c^0(1 - e^{-\alpha\bar{T}_L}) = \frac{c^0}{f_1(\alpha\bar{T}_L)}, \\ f_1(\alpha t^-) &\geq f_1(\alpha\bar{T}_L). \end{aligned}$$

As  $f_1(t)$  is a decreasing function of  $t$ , we have  $t^- \leq \bar{T}_L$ . Since

$$\left( \min_{r \in R} d_r \right) \min_{i \in N_0} \left( \frac{h_i}{\alpha} + u_i \right) \frac{\bar{T}_L}{1 - e^{-\alpha\bar{T}_L}} \leq \frac{H(N)\bar{T}_L}{1 - e^{-\alpha\bar{T}_L}} \leq c^0 = \min_{r \in R} d_r \min_{i \in N_0} \left( \frac{h_i}{\alpha} + u_i \right) \frac{t^+}{1 - e^{-\alpha t^+}}.$$

Therefore,

$$\frac{\bar{T}_L}{1 - e^{-\alpha\bar{T}_L}} \leq \frac{t^+}{1 - e^{-\alpha t^+}},$$

or  $f_2(\bar{T}_L) \leq f_2(\alpha t^+)$ . As  $f_2(t)$  is an increasing function of  $t$ , we have  $\bar{T}_L \leq t^+$ .

Therefore, vector  $\bar{T}_L 1 \in \theta$  and we have

$$TCPV(t^*, k^*) \leq c^0. \quad (\text{B.1})$$

In view of (B.1), it suffices to show that

$$TCPV(t, k) \geq c^0 \text{ for all } t \in T \setminus \theta \text{ and } k \in K. \quad (\text{B.2})$$

Thus, fix  $t \in T \setminus \theta$  and  $k \in K$ . There are two cases to consider:

$$(1) t_a = \min_{i \in N} t_i < t^- \text{ for some } a \in N.$$

Clearly,

$$TCPV(t, k) \geq k_a f_1(\alpha t_a) \geq \delta f_1(\alpha t_a) \geq \delta f_1(\alpha t^-) = c^0. \quad (\text{B.3})$$

$$(2) \max_{r \in R} \tau_r > t^+.$$

Note, in view of (2.7), that  $\max \tau_r$  is achieved for some route  $\ell$  whose initial node  $a \in N_0$ . Hence,

$$TCPV(t, k) \geq \sum_{r \in R} \left( u_{i_r} + \frac{h_{i_r}}{\alpha} \right) \frac{d_r}{\alpha} f_2(\alpha \tau_r) \geq H_\ell f_2(t^+) \quad (\text{B.4})$$

$$\geq \left( \min_{r \in R} d_r \right) \min_{i \in N_0} \left( \frac{h_i}{\alpha} + u_i \right) \alpha / f_2(t^+) = c^0. \quad (\text{B.5})$$

Thus (B.2) follows from (B.3) and (B.5).  $\square$

### Proof of Lemma 3.3 (Another expression of the continuous relaxation (RP)):

For any  $r \in R$  and  $i \in R$ . let  $x_{ri}$  be a Lagrange multiplier associated with constraint (2.8). Fix  $k^0 \in K$ , let  $(RP_{k^0})$  denote the convex (inner) minimization problem to the right of (3.2) with  $k_i = k_i^0 (i \in N)$  and let  $TCPV^*(k^0)$  denote its optimal value. That is, the problem  $(RP_{k^0})$  is

$$TCPV^*(k^0) = \min f(t, k^0),$$

subject to

$$f_2(\alpha \tau_r) \geq f_2(\alpha t_i), \quad \forall r \in R, i \in r,$$

$$\tau, t > 0.$$

Let  $f(t, k^0) = \alpha^2/2 \times TCPV(t, k^0)$ , where  $TCPV(t, k^0)$  is defined in (2.6). The associated Lagrangian of  $(RP_{k^0})$  is given by

$$L(x, t) = \sum_{i \in N} \frac{k_i^0 \alpha^2}{2} f_1(\alpha t_i) + \sum_{r \in R} H_r f_2(\alpha \tau_r) - \frac{1}{2} \sum_{r \in R} h_{i_r} d_r + \sum_{r \in R} \sum_{i \in r} x_{ri} [f_2(\alpha t_i) - f_2(\alpha \tau_r)].$$

By definition, see Geoffrion [9] (1971), the Lagrangian dual of  $(RP_{k^0})$  can be written as

$$D = \sup_{x \geq 0} \inf_{t > 0} L(x, t). \quad (\text{B.6})$$

Regrouping the terms of  $L(x, t)$ , we have

$$L(x, t) = \sum_{i \in N} \left[ \frac{k_i^0 \alpha^2}{2} f_1(\alpha t_i) + \left( \sum_{r: i \in r} x_{ri} \right) f_2(\alpha t_i) \right] + \sum_{r \in R} \left[ H_r - \sum_{i \in r} x_{ri} \right] f_2(\alpha \tau_r) - \frac{1}{2} \sum_{r \in R} h_{i_r} d_r. \quad (\text{B.7})$$

If, for a given  $x \geq 0$ ,  $\left( H_r - \sum_{i \in r} x_{ri} \right) < 0$ , then  $\inf_{t > 0} L(x, t) = \lim_{t \rightarrow \infty} L(x, t) = -\infty$ ; likewise, if  $\left( H_r - \sum_{i \in r} x_{ri} \right) \geq 0$ ,  $\inf_{\tau_r > 0} \left[ H_r - \sum_{i \in r} x_{ri} \right] f_2(\alpha \tau_r) = 0$ . To achieve the supremum in (B.6), we may thus restrict ourselves to vectors  $x$  for which

$$H_r \geq \sum_{i \in r} x_{ri}, \quad r \in R. \quad (\text{B.8})$$

Note that the last summation  $\frac{1}{2} \sum_{r \in R} h_{i_r} d_r$  in (B.7) is a constant. We will allocate it to individual item  $i$  by  $w_i$  as defined in (3.3f). Note also that in an optimal solution  $(x^*, v^*)$ , (B.8) must be satisfied as *equalities*, we have

$$\sum_{r \in R} \frac{h_{i_r} d_r}{2} = \sum_{r \in R} (H_r y_r) = \sum_{r \in R} \left( \sum_{i \in r} x_{ri} y_r \right) = \sum_{i \in N} \left( \sum_{r: i \in r} x_{ri} y_r \right) = \sum_{i \in N} w_i.$$

Also we note that this allocation will not affect the optimal solution. However,  $w_i$  has to satisfy the inequality (3.3f) in order to use the estimation of the inventory cost for power-of-two policies. Let

$$y_i \triangleq \max_{r: i \in r} y_r, \quad \forall i \in N.$$

Obviously,  $y_r < 1$ ,  $y_i < 1$ , and we can see that  $w_i$  satisfying the inequality (3.3f) due to the following inequality:

$$w_i \leq \sum_{r: i \in r} x_{ri} y_i = v_i y_i < v_i.$$

Imposing these additional constraints and substituting (3.3c), we may rewrite (B.6) as:

$$\begin{aligned}
 D &= \sup_{\{x \geq 0, v \geq 0: (3.3c), (B.8)\}} \inf_{t > 0} \sum_{i \in N} \left[ k_i^0 (\alpha^2/2) f_1(\alpha t_i) + v_i f_2(\alpha t_i) - w_i \right] \\
 &= \sup_{\{x \geq 0, v \geq 0: (3.3c), (B.8)\}} \sum_{i \in N} \min_{t > 0} \left[ \frac{k_i^0 \alpha^2/2 + v_i \alpha t_i}{1 - e^{-\alpha t_i}} - w_i \right] \\
 &= \sup_{\{x \geq 0, v \geq 0: (3.3c), (B.8)\}} \frac{\alpha^2}{2} \sum_{i \in N} TCPV(\alpha t_i^* | k_i^0, v_i, w_i), \tag{B.9}
 \end{aligned}$$

where  $TCPV(\alpha t_i^* | k_i^0, v_i, w_i)$  is defined in (1.6). *Strong duality*, i.e.,  $TCPV^*(k^0) = D$ , and the existence of a pair  $(x^*, v^*)$  achieving the supremum in (B.9) all follow from  $(RP_k)$  being a *stable convex* program, see Theorem 3 in Geoffrion [9] (1971). Stability may be verified by Slater's condition, i.e., there exists a vector  $t > 0$ , with  $\tau_r > T_i$  for all  $r \in R$  and  $i \in r$ . (Let  $\tau_r = 2$ , all  $r \in R$  and  $T_i = 1$ , all  $i \in N$ .) Finally note that in an optimal solution  $(x^*, v^*)$ , (B.8) must be satisfied as *equalities* and may thus be replaced by (3.3b). The lemma thus follows from (B.9) and the subsequent observations.  $\square$

## C Proofs of Theorems and Lemmas in section 4

### Proof of Theorem 4.1 (The Lower Bound Theorem):

Assume first that  $K(\cdot)$  is strictly monotone so that Lemmas 3.2 and 3.3 apply. Let  $t^*$  achieve the minimum in  $(RP)$ , see Lemma 3.2. It follows from Lemma 3.3 that a vector  $k^* \in K$ , and vectors  $x^*$ ,  $v^*$  exist which satisfy (3.3b), (3.3c), (3.3d) and with  $TCPV^* = \sum_{j \in N} TCPV(\alpha t_j^* | k^*, v^*, w_j)$ .

For a feasible policy  $\pi$ , let  $c$  be the NPV of the total cost incurred by the policy. We show that  $c \geq TCPV^*$ . We evaluate the total setup costs, production cost and holding costs separately.

Let  $\eta^m$  ( $m = 0, 1, 2, \dots$ ) be all the instants at which items are replenished, and  $S_m$  be the set of items replenished at instant  $\eta^m$ . Let  $\eta_j^\ell$  (or  $0 = \eta_j^0, \eta_j^1, \eta_j^2, \dots$ ) be the instant at which item  $j$  is replenished at its  $\ell$ th time and the order quantity for item  $j$  is  $Q_j^\ell$ . Let  $\tau_r^\ell$  (or  $0 = \tau_r^0, \tau_r^1, \tau_r^2, \dots$ ) be the instant at which echelon inventory on route  $r$  is replenished at its  $\ell$ th time.

Note that  $\eta_j^\ell$  is simply the renumbering of the replenish instant of  $\eta^m$  in terms of item  $j$  in order to make the summation over product  $j$  easier to express.

The NPV of the total setup cost is

$$\begin{aligned} K^{tot} &= \sum_{m=0}^{\infty} K(S_m) e^{-\alpha \eta^m} \geq \sum_{m=0}^{\infty} \left( \sum_{j \in S_m} k_j \right) e^{-\alpha \eta^m} = \sum_{j \in N} \left( \sum_{\ell=0}^{\infty} k_j e^{-\alpha \eta_j^\ell} \right) \\ &= \sum_{j \in N} k_j \left( \sum_{\ell=0}^{\infty} e^{-\alpha \eta_j^\ell} \right). \end{aligned}$$

Note that the first node of route  $r$  is  $i_r$ . The holding cost on route  $r$  is:

$$\begin{aligned} &\sum_{\ell=0}^{\infty} \left( h_{i_r} d_r \int_0^{\eta_r^{\ell+1} - \eta_r^\ell} (\eta_r^{\ell+1} - \eta_r^\ell - t) e^{-\alpha t} dt \right) e^{-\alpha \eta_r^\ell} \\ &= \frac{h_{i_r} d_r}{\alpha} \sum_{\ell=0}^{\infty} \left( \eta_r^{\ell+1} - \eta_r^\ell - \frac{1 - e^{-\alpha(\eta_r^{\ell+1} - \eta_r^\ell)}}{\alpha} \right) e^{-\alpha \eta_r^\ell} \\ &= \frac{h_{i_r} d_r}{\alpha} \sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^\ell) e^{-\alpha \eta_r^\ell} - \frac{h_{i_r} d_r}{\alpha^2} \sum_{\ell=0}^{\infty} (e^{-\alpha \eta_r^\ell} - e^{-\alpha \eta_r^{\ell+1}}) \\ &= \frac{h_{i_r} d_r}{\alpha} \sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^\ell) e^{-\alpha \eta_r^\ell} - \frac{h_{i_r} d_r}{\alpha^2}. \end{aligned}$$

The NPV of the total holding cost is

$$H^{tot} = \sum_{r \in R} \left[ \frac{h_{i_r} d_r}{\alpha} \sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^\ell) e^{-\alpha \eta_r^\ell} \right] - \sum_{r \in R} \frac{h_{i_r} d_r}{\alpha^2}.$$

The NPV of the total production cost is

$$U^{tot} = \sum_{r \in R} \left[ u_{i_r} d_r \sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^\ell) e^{-\alpha \eta_r^\ell} \right].$$

And the NPV of the total holding cost and production cost is

$$H^{tot} + U^{tot} = \sum_{r \in R} \left[ \left( \frac{h_{i_r}}{\alpha} + u_{i_r} \right) \frac{d_r}{\alpha} \sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^\ell) \alpha e^{-\alpha \eta_r^\ell} \right] - \sum_{r \in R} \frac{h_{i_r} d_r}{\alpha^2}.$$

If  $t_3 > t_2 > t_1 \geq 0$ , then it is obvious that

$$(t_3 - t_1) e^{-\alpha t_1} \geq (t_2 - t_1) e^{-\alpha t_1} + (t_3 - t_2) e^{-\alpha t_2}.$$

This means whenever we add extra replenishment instants, the summation  $\sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^\ell) e^{-\alpha \eta_r^\ell}$  is getting smaller. Therefore, for any node  $i \in r$  the replenishment set  $(\eta_r^0, \eta_r^1, \eta_r^2, \dots)$  of route  $r$  is the

subset of the replenishment set  $(\eta_i^0, \eta_i^1, \eta_i^2, \dots)$  of node  $i$ . And we have

$$\sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^{\ell}) e^{-\alpha\eta_r^{\ell}} \geq \sum_{\ell=0}^{\infty} (\eta_i^{\ell+1} - \eta_i^{\ell}) e^{-\alpha\eta_i^{\ell}}, \quad \forall i \in r.$$

Therefore, we can get the lower bound of the NPV of the total holding cost and production cost

$$\begin{aligned} H^{tot} + U^{tot} &= \sum_{r \in R} \left[ \left( \frac{h_{i_r}}{\alpha} + u_{i_r} \right) \frac{d_r}{\alpha} \sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^{\ell}) \alpha e^{-\alpha\eta_r^{\ell}} \right] - \sum_{r \in R} \frac{h_{i_r} d_r}{\alpha^2} \\ &= \sum_{r \in R} \left[ \left( \sum_{i \in r} x_{ri}^* \right) \sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^{\ell}) \alpha e^{-\alpha\eta_r^{\ell}} \right] - \sum_{r \in R} \frac{h_{i_r} d_r}{\alpha^2} \\ &= \sum_{i \in N} \sum_{r: i \in r} \left\{ x_{ri}^* \left[ \sum_{\ell=0}^{\infty} (\eta_r^{\ell+1} - \eta_r^{\ell}) \alpha e^{-\alpha\eta_r^{\ell}} \right] \right\} - \sum_{r \in R} \frac{h_{i_r} d_r}{\alpha^2} \\ &\geq \sum_{i \in N} \sum_{r: i \in r} \left\{ x_{ri}^* \left[ \sum_{\ell=0}^{\infty} (\eta_i^{\ell+1} - \eta_i^{\ell}) \alpha e^{-\alpha\eta_i^{\ell}} \right] \right\} - \sum_{r \in R} \frac{h_{i_r} d_r}{\alpha^2} \\ &= \sum_{i \in N} \left\{ \left( \sum_{r: i \in r} x_{ri}^* \right) \left[ \sum_{\ell=0}^{\infty} (\eta_i^{\ell+1} - \eta_i^{\ell}) \alpha e^{-\alpha\eta_i^{\ell}} \right] \right\} - \sum_{r \in R} \frac{h_{i_r} d_r}{\alpha^2} \\ &= \sum_{i \in N} \left\{ v_i^* \sum_{\ell=0}^{\infty} (\eta_i^{\ell+1} - \eta_i^{\ell}) \alpha e^{-\alpha\eta_i^{\ell}} - w_i \right\}. \end{aligned}$$

The NPV  $c$  of the total cost is therefore,

$$\begin{aligned} c &= K^{tot} + H^{tot} + U^{tot} \geq \sum_{i \in N} \left\{ k_i \left( \sum_{\ell=0}^{\infty} e^{-\alpha\eta_i^{\ell}} \right) + v_i^* \sum_{\ell=0}^{\infty} (\eta_i^{\ell+1} - \eta_i^{\ell}) \alpha e^{-\alpha\eta_i^{\ell}} - w_i \right\} \quad (C.1) \\ &\geq \sum_{i \in N} TCPV(\alpha t_i^* | k_i, v_i^*, w_i) = TCPV^*. \end{aligned}$$

The last inequality is based on that  $TCPV(\alpha t_i^* | k_i, v_i^*, w_i)$  is the optimal solution for each term in (C.1) satisfying (3.3f), which is concluded from the result on the single product expressed in (1.2) and satisfies (1.3b). Therefore, the optimal value  $TCPV^*$  of  $(RP)$  is a lower bound on the total cost (NPV) for any feasible policy.  $\square$

### Proof of Lemma 4.2:

Let

$$f(\beta) = \frac{1}{\beta \ln 2},$$

then we have

$$\int_{\sqrt{0.5}}^{\sqrt{2}} f(\beta) d\beta = 1.$$

Based on equation (4.1), we have

$$\frac{\bar{t}_i}{t_i} = \begin{cases} \frac{2^{p_i} \beta T_0}{b_i T_0 2^{p_i} \sqrt{0.5}} = \frac{\beta \sqrt{2}}{b_i}, & \text{if } \sqrt{0.5} \leq \beta < b_i, \\ \frac{2^{p_i-1} \beta T_0}{b_i T_0 2^{p_i} \sqrt{0.5}} = \frac{\beta \sqrt{0.5}}{b_i}, & \text{if } b_i \leq \beta < \sqrt{2}. \end{cases}$$

And therefore,

$$\begin{aligned} \int_{\sqrt{0.5}}^{\sqrt{2}} TC(\bar{t}_i) \frac{d\beta}{\beta \ln 2} &= \int_{\sqrt{0.5}}^{\sqrt{2}} TC(\bar{t}_i) f(\beta) d\beta \\ &= \int_{\sqrt{0.5}}^{b_i} TC\left(\frac{\beta \sqrt{2}}{b_i} t_i\right) f(\beta) d\beta + \int_{b_i}^{\sqrt{2}} TC\left(\frac{\beta}{b_i \sqrt{2}} t_i\right) f(\beta) d\beta \end{aligned} \quad (C.2)$$

$$\begin{aligned} &= \int_{\frac{1}{b_i}}^{\sqrt{2}} TC(rt_i) f\left(\frac{rb_i}{\sqrt{2}}\right) \frac{b_i}{\sqrt{2}} dr + \int_{\sqrt{0.5}}^{\frac{1}{b_i}} TC(rt_i) f(rb_i \sqrt{2}) b_i \sqrt{2} dr \\ &= \int_{\frac{1}{b_i}}^{\sqrt{2}} TC(rt_i) \frac{\sqrt{2}}{rb_i \ln 2} \frac{b_i}{\sqrt{2}} dr + \int_{\sqrt{0.5}}^{\frac{1}{b_i}} TC(rt_i) \frac{1}{rb_i \sqrt{2} \ln 2} b_i \sqrt{2} dr \\ &= \int_{\sqrt{0.5}}^{\sqrt{2}} TC(rt_i) \frac{dr}{r \ln 2}. \end{aligned} \quad (C.3)$$

Note that from (C.2) to (C.3) we make a replacement  $r = \frac{\beta \sqrt{2}}{b_i}$  for the first integration while we make a replacement  $r = \frac{\beta}{b_i \sqrt{2}}$  for the second integration. As we have

$$TC(rt_i) \leq \frac{1+\epsilon}{2} \left(r + \frac{1}{r}\right) TC(t_i),$$

therefore,

$$\begin{aligned} \int_{\sqrt{0.5}}^{\sqrt{2}} TC(rt_i) \frac{dr}{r \ln 2} &\leq \int_{\sqrt{0.5}}^{\sqrt{2}} \frac{1+\epsilon}{2} \left(r + \frac{1}{r}\right) TC(t_i) \frac{dr}{r \ln 2} \\ &= TC(t_i) \frac{1+\epsilon}{2} \int_{\sqrt{0.5}}^{\sqrt{2}} \left(r + \frac{1}{r}\right) \frac{dr}{r \ln 2} \\ &= \frac{1+\epsilon}{\sqrt{2} \ln 2} TC(t_i). \end{aligned}$$

This completes the proof. □

### Proof of Theorem 4.3 (Worst case analysis of power-of-two policies):

We assume that  $K(\cdot)$  is strictly monotone so that Lemma 3.3 applies. If it is not, the proof can be amended by considering a sequence of perturbed setup cost functions. In view of Lemma 3.3, let  $t^* = \{t_1^*, t_2^*, \dots, t_n^*\}$  denote an optimal solution to  $(RP)$  and  $\bar{t}^* = (\bar{t}_1^*, \bar{t}_2^*, \dots, \bar{t}_n^*)$  the power-of-two

vector obtained by the rounding procedure. Let  $\gamma = (\gamma_1, \gamma_2, \dots, \gamma_n)$  be a permutation of the node indices such that  $t_{\gamma_1}^{*-1} \geq t_{\gamma_2}^{*-1} \geq \dots \geq t_{\gamma_n}^{*-1}$ . Note that the components of  $\bar{t}^*$  may be ranked in the same way.

Based on Lemma 1.3,  $\alpha t_i^*$  is the optimal solution to  $TCPV(\alpha t_i^* | k_i, v_i, w_i)$ . And  $\bar{t}_i^* = r_i t_i^*$  the power-of-two rounding of  $t_i^*$  with  $r_i \in (\sqrt{0.5}, \sqrt{2}]$  for all  $i \in N$ . Based on Lemma 1.3, we have

$$\frac{TCPV(r_i \alpha t_i^* | k_i, v_i, w_i)}{TCPV(\alpha t_i^* | k_i, v_i, w_i)} \leq \frac{1 + \epsilon}{2} \left( r_i + \frac{1}{r_i} \right), \quad \forall r_i \in [\sqrt{0.5}, \sqrt{2}], \quad i \in N,$$

where  $\epsilon = 0.000570126 > 0$ .

It is easy to verify that

$$r + \frac{1}{r} \leq \frac{3}{\sqrt{2}} = 2.121320344, \quad \forall r \in [\sqrt{0.5}, \sqrt{2}].$$

And we have the following estimate for fixed base period  $T_L$ :

$$\begin{aligned} \sum_{i \in N} TCPV(\alpha \bar{t}_i^* | k_i, v_i, w_i) &= \sum_{i \in N} TCPV(r_i \alpha t_i^* | k_i, v_i, w_i) \\ &\leq \sum_{i \in N} \frac{1 + \epsilon}{2} \left( r_i + \frac{1}{r_i} \right) TCPV(\alpha t_i^* | k_i, v_i, w_i) \\ &\leq (1 + \epsilon) \frac{2.121320344}{2} \sum_{i \in N} TCPV(\alpha t_i^* | k_i, v_i, w_i) \\ &\leq 1.061264882 \times TCPV^*. \end{aligned}$$

Based on Lemma 4.2, we have the following estimate for variable base period  $T_L$ :

$$\begin{aligned} \min_{T_L > 0} \sum_{i \in N} TCPV(\alpha \bar{t}_i^* | k_i, v_i, w_i) &= \min_{T_L > 0} \sum_{i \in N} TCPV(r_i \alpha t_i^* | k_i, v_i, w_i) \\ &\leq \sum_{i \in N} \int_{\sqrt{0.5}}^{\sqrt{2}} TCPV(\alpha t_i^* | k_i, v_i, w_i) \frac{dr_i}{r \ln 2} \\ &\leq \sum_{i \in N} \frac{1 + \epsilon}{\sqrt{2} \ln 2} TCPV(\alpha t_i^* | k_i, v_i, w_i) \\ &\leq \frac{1 + \epsilon}{\sqrt{2} \ln 2} \sum_{i \in N} TCPV(\alpha t_i^* | k_i, v_i, w_i) \\ &\leq 1.020721055 \times TCPV^*. \end{aligned}$$

This completes our proof. □

## D Proofs of Theorems and Lemmas in section 5

### Proof of Lemma 5.1

Let  $T$  be the reorder interval. The total average cost  $TC$  for EOQ model is:

$$TC = \frac{k}{T} + \frac{1}{2}(h + \alpha u)dT = \frac{k}{T} + vdT,$$

where  $v$  is defined in (1.3a). We can easily find the optimal reorder interval  $T^{EOQ}$  below.

$$T^{EOQ} = \sqrt{k/v}. \quad (D.1)$$

If the reorder interval is  $T$ , the NPV of the total cost  $TCPV(\alpha T|k, v, w)$  is defined in (1.5).

$$TCPV(\alpha T|k, v, w) = \frac{2}{\alpha^2} \left( \frac{k(\alpha^2/2) + v\alpha T}{1 - e^{-\alpha T}} - w \right),$$

where  $w$  is defined in (1.3b). Then

$$\alpha T^{EOQ} = \sqrt{2q},$$

where  $q$  is defined in (1.4d). We have

$$\begin{aligned} TCPV(\alpha T^{EOQ}|k, v, w) &= TCPV(\sqrt{2q}|k, v, w) = \frac{2v}{\alpha^2} \left( \frac{(k/v)(\alpha^2/2) + \sqrt{2q}}{1 - e^{-\sqrt{2q}}} - \frac{w}{v} \right) \\ &= \frac{2v}{\alpha^2} \left( \frac{q + \sqrt{2q}}{1 - e^{-\sqrt{2q}}} - \frac{w}{v} \right). \end{aligned}$$

The optimal solution  $\theta^*$  to  $TCPV(\theta|k, v, w)$  and the optimal value of  $TCPV(\theta|k, v, w)$  can be found in Lemma 1.2. And the optimal reorder interval is  $T^{NPV} = \theta^*/\alpha$ .

$$TCPV(\theta^*|k, v, w) = TCPV(\alpha T^{NPV}|k, v, w) = \frac{2v}{\alpha^2} \left( -LambertW(-1, -e^{-1-q}) - \frac{w}{v} \right).$$

Therefore, the ratio of  $TCPVs$  based on EOQ model and NPV model:

$$\begin{aligned} \frac{TCPV(\alpha T^{EOQ}|k, v, w)}{TCPV(\alpha T^{NPV}|k, v, w)} &= \frac{\frac{2v}{\alpha^2} \left( \frac{q + \sqrt{2q}}{1 - e^{-\sqrt{2q}}} - \frac{w}{v} \right)}{\frac{2v}{\alpha^2} \left( -LambertW(-1, -e^{-1-q}) - \frac{w}{v} \right)} \\ &\leq \frac{\frac{q + \sqrt{2q}}{1 - e^{-\sqrt{2q}}} - 1}{-LambertW(-1, -e^{-1-q}) - 1} < 1.09603. \end{aligned}$$

where  $q^* = 28.75$  gives us the largest value of the ratio above and is the base for the last inequality of  $q > 0$ . □

### Proof of Lemma 5.4

1. We have

$$\begin{aligned}
 c^{EOQ} &= \max_{i \in N} \sum 2(k_i v_i^*)^{1/2}, \\
 H_r &= \sum_{i \in r} x_{ri}^*, \quad r \in R, \\
 \sum_{r: i \in r} x_{ri}^* &= v_i^*, \quad i \in N, \\
 k &\geq 0, x \geq 0, v \geq 0.
 \end{aligned} \tag{D.2}$$

Note that  $\tau_r \geq t_i$  for all  $i \in r$ . We have

$$\sum_{r \in R} H_r \tau_r = \sum_{r \in R} \sum_{i \in r} x_{ri}^* \tau_r = \sum_{i \in N} \sum_{r: i \in r} (x_{ri}^* \tau_r) \geq \sum_{i \in N} \sum_{r: i \in r} x_{ri}^* t_i = \sum_{i \in N} v_i^* t_i.$$

Therefore,

$$c^{EOQ} = \min_{t \in T} \max_{k \in K} \left\{ \sum_{i \in N} k_i / t_i + \sum_{r \in R} H_r \tau_r \right\} \geq \min_{t \in T} \sum_{i \in N} \{k_i^* / t_i + v_i^* t_i\}.$$

Let

$$\begin{aligned}
 t_i^{EOQ} &\triangleq \sqrt{k_i^* / v_i^*} = \frac{1}{\alpha} \sqrt{2q_i^*}, \\
 \tau_r^{EOQ} &\triangleq \max_{i \in r} t_i^{EOQ}, \quad \forall r \in R,
 \end{aligned}$$

which minimize the equation above and  $q_i^*$  is defined in (5.1). Combining the equation (D.2),

we have

$$c^{EOQ} \geq \min_{t \in T} \sum_{i \in N} \{k_i^* / t_i + v_i^* t_i\} = \sum_{i \in N} 2(k_i^* v_i^*)^{1/2} = c^{EOQ}.$$

This means we can replace the equality for all the inequalities above. Then we have

$$\sum_{r:i \in r} (x_{ri}^* \tau_r^{EOQ}) = \sum_{r:i \in r} x_{ri}^* t_i^{EOQ}, \quad \forall i \in N.$$

As  $\tau_r^{EOQ} \geq t_i^{EOQ}$  for all  $i \in r$  and  $x_{ri}^* \geq 0$ , we conclude that

$$\text{If } \tau_r^{EOQ} > t_i^{EOQ}, \text{ then } x_{ri}^* = 0, \quad (\text{D.3a})$$

$$\text{If } x_{ri}^* > 0, \text{ then } \tau_r^{EOQ} = t_i^{EOQ}. \quad (\text{D.3b})$$

As

$$\begin{aligned} \sum_{r \in R} H_r \tau_r^{EOQ} &= \sum_{r \in R} \sum_{i \in r} x_{ri}^* \tau_r^{EOQ} = \sum_{i \in N} \sum_{r:i \in r} x_{ri}^* \tau_r^{EOQ} = \sum_{i \in N} \sum_{\substack{r:i \in r \\ x_{ri}^* > 0}} x_{ri}^* \tau_r^{EOQ} \\ &= \sum_{i \in N} \sum_{\substack{r:i \in r \\ x_{ri}^* > 0}} x_{ri}^* t_i^{EOQ} = \sum_{i \in N} \sum_{r:i \in r} x_{ri}^* t_i^{EOQ} = \sum_{i \in N} v_i^* t_i^{EOQ}. \end{aligned}$$

Therefore,

$$c^{EOQ} = \sum_{i \in N} k_i^* / t_i^{EOQ} + \sum_{r \in R} H_r \tau_r^{EOQ} = \sum_{i \in N} (k_i^* / t_i^{EOQ} + v_i^* t_i^{EOQ}) = \sum_{i \in N} 2(k_i^* v_i^*)^{1/2}.$$

That is,  $t_i^{EOQ}$  and  $\tau_r^{EOQ}$  are the optimal solution to the average cost model. Now we would like to calculate the NPV of total cost in the NPV model using  $t_i^{EOQ}$  and  $\tau_r^{EOQ}$ . Based on equation (2.6), the lower bound on the NPV of the total cost is as follows.

$$TCPV(t, k) = \frac{2}{\alpha^2} \left\{ \sum_{i \in N} k_i^* (\alpha^2 / 2) f_1(\alpha t_i^{EOQ}) + \sum_{r \in R} H_r f_2(\alpha \tau_r^{EOQ}) - \sum_{r \in R} \frac{h_r d_r}{2} \right\}.$$

Because

$$\begin{aligned} \sum_{r \in R} H_r f_2(\alpha \tau_r^{EOQ}) &= \sum_{r \in R} \sum_{i \in r} x_{ri}^* f_2(\alpha \tau_r^{EOQ}) = \sum_{i \in N} \sum_{r:i \in r} x_{ri}^* f_2(\alpha \tau_r^{EOQ}) \\ &= \sum_{i \in N} \sum_{\substack{r:i \in r \\ x_{ri}^* > 0}} x_{ri}^* f_2(\alpha \tau_r^{EOQ}) = \sum_{i \in N} \sum_{\substack{r:i \in r \\ x_{ri}^* > 0}} x_{ri}^* f_2(\alpha t_i^{EOQ}) \\ &= \sum_{i \in N} \left[ \sum_{r:i \in r} x_{ri}^* \right] f_2(\alpha t_i^{EOQ}) = \sum_{i \in N} v_i^* f_2(\alpha t_i^{EOQ}), \end{aligned}$$

and let It is easy to verify that

$$\sum_{r \in R} \frac{h_{i_r} d_r}{2} = \sum_{i \in N} w_i^*.$$

Therefore, we have

$$\begin{aligned} TCPV(t^{EOQ}, k) &= \frac{2}{\alpha^2} \sum_{i \in N} \left\{ k_i^* (\alpha^2/2) f_1(\alpha t_i^{EOQ}) + v_i^* f_2(t_i^{EOQ}) - w_i^* \right\} \\ &= \sum_{i \in N} \frac{2v_i^*}{\alpha^2} \left[ \frac{(k_i^*/v_i^*)(\alpha^2/2) + \alpha t_i^{EOQ}}{1 - e^{-\alpha t_i^{EOQ}}} - \frac{w_i^*}{v_i^*} \right] \\ &= \sum_{i \in N} \frac{2v_i^*}{\alpha^2} \left[ \frac{q_i^* + \sqrt{2q_i^*}}{1 - e^{-\sqrt{2q_i^*}}} - \frac{w_i^*}{v_i^*} \right]. \end{aligned}$$

2. Now we will find a feasible solution to NPV model based on  $x_{ri}^*$  and  $v_i^*$ .

Based on Lemma 1.2, we have

$$\begin{aligned} t_i^{NPV} &= \frac{1}{\alpha} [-1 - q_i^* - LambertW(-1, -e^{-1-q_i^*})], \\ TCPV(\alpha t_i^{NPV} | k, v_i^*, w_i^*) &= \frac{2v_i^*}{\alpha^2} \left[ -LambertW(-1, -e^{-1-q_i^*}) - \frac{w_i^*}{v_i^*} \right], \\ \tau_r^{NPV} &\triangleq \max_{i \in r} t_i^{NPV}. \end{aligned}$$

Because both  $t_i^{EOQ}$  and  $t_i^{NPV}$  are strictly increasing functions of  $q_i^*$ , the relationships in (D.3) will hold for  $t_i^{NPV}$  also. That is,

$$\text{If } \tau_r^{NPV} > t_i^{NPV}, \text{ then } x_{ri}^* = 0, \quad (\text{D.4a})$$

$$\text{If } x_{ri}^* > 0, \text{ then } \tau_r^{NPV} = t_i^{NPV}. \quad (\text{D.4b})$$

We know that  $k_i^*, t_i^{NPV}, \tau_r^{NPV}$  are a feasible solution to continuous relaxation (RP). Therefore, based on the first equality in Lemma 3.2, we have

$$TCPV(t^{NPV}, k) \geq TCPV^*. \quad (\text{D.5})$$

However, we also have

$$\begin{aligned}
 \sum_{r \in R} H_r f_2(\alpha \tau_r^{NPV}) &= \sum_{r \in R} \sum_{i \in r} x_{ri}^* f_2(\alpha \tau_r^{NPV}) = \sum_{i \in N} \sum_{r: i \in r} x_{ri}^* f_2(\alpha \tau_r^{NPV}) \\
 &= \sum_{i \in N} \sum_{\substack{r: i \in r \\ x_{ri}^* > 0}} x_{ri}^* f_2(\alpha \tau_r^{NPV}) = \sum_{i \in N} \sum_{\substack{r: i \in r \\ x_{ri}^* > 0}} x_{ri}^* f_2(\alpha t_i^{NPV}) \\
 &= \sum_{i \in N} \left[ \sum_{r: i \in r} x_{ri}^* \right] f_2(\alpha t_i^{NPV}) = \sum_{i \in N} v_i^* f_2(\alpha t_i^{NPV}).
 \end{aligned}$$

Therefore,

$$\begin{aligned}
 TCPV(t^{NPV}, k) &= \frac{2}{\alpha^2} \left\{ \sum_{i \in N} k_i^* (\alpha^2/2) f_1(\alpha t_i^{NPV}) + \sum_{r \in R} H_r f_2(\alpha \tau_r^{NPV}) - \sum_{r \in R} \frac{h_{i_r} d_r}{2} \right\} \\
 &= \sum_{i \in N} \frac{2}{\alpha^2} \left\{ k_i^* (\alpha^2/2) f_1(\alpha t_i^{NPV}) + v_i^* f_2(\alpha t_i^{NPV}) - w_i^* \right\} \\
 &= \sum_{i \in N} \frac{2v_i^*}{\alpha^2} \left[ \frac{q_i^* + \alpha t_i^{NPV}}{1 - e^{\alpha t_i^{NPV}}} - \frac{w_i^*}{v_i^*} \right] \\
 &= \sum_{i \in N} \frac{2v_i^*}{\alpha^2} \left[ -\text{LambertW}(-1, -e^{-1-q_i^*}) - \frac{w_i^*}{v_i^*} \right].
 \end{aligned}$$

That is,  $k_i^*$ ,  $x_i^*$ ,  $v_i^*$  are also a feasible solution to the continuous relaxation (RP) in Lemma 3.3.

That is,

$$TCPV(t^{NPV}, k) \leq TCPV^*. \tag{D.6}$$

From (D.5) and (D.6), we conclude that

$$TCPV(t^{NPV}, k) = TCPV^*.$$

This completes our proof. □